REST Journal on Emerging trends in Modelling and Manufacturing
Vol:4(4),2018
REST Publisher
ISSN: 2455-4537

Website: www.restpublisher.com/journals/jemm

Total Domination in Graphs: New Perspectives and Results

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Abstract

In this research paper, we initiate a study of such questions in the present context of total domination theory within graph theory from new perspectives and results concerned with formal definitions, mathematical formulations, theorems and practical applications. Description Total domination is a tool used to determine the graph that has minimum number of vertices can monitor or control all other vertices - factors of interest for problems related, amongst others fields as network optimization and resilience analysis. This paper covers fundamental definitions as well their most recent theoretical bounds, algorithmic complexities and practical examples arising in applications from the real world. The obtained results could be helpful in deciphering, exploiting and administering the concept of total domination further.

Keywords: Total Domination, Graph Theory, Mathematical Formulations, Theoretical Bounds, Algorithmic Complexity, Network Optimization, Resilience Analysis.

1. Introduction

Background on Total Domination in Graph Theory

In the study of graph theory there is an important notion called Total Domination that deals with finding (size) smallest subset on a given set of vertices V in $(G = (V, E)$, such every vertex either belongs to this subset or has at least one adjacent neighbour node within the mean subset [1]. A set $D \subseteq V$ is a total dominating if it satisfies the following formal definition:

$$\forall \forall v \in V \setminus D, \exists u \in D \text{ such that } \{u, v\} \in E$$

The total domination number represented as $\gamma_t(G)$of G is the property of cardinality of the smallest total dominating set D. This concept provides a measure of the efficiency with which a subset of vertices can control or influence the entire graph [1].

Importance and Applications of Total Domination

All or nothing dominates in all the practical and theoretical applications. Network theory use is valuable for understanding and improving communication networks in which desired nodes (vertices) should sense or control other ones. It is also widely to locate facilities, the task of selecting a minimum number of locations (vertices) that in turns span all facilities(vertices) [2].

Overview of Traditional Results and Motivations for New Perspectives

Previous research on total domination has generally been dedicated to developing efficient algorithms for computing $\gamma_t(G)$, as well as bounding its value, and relating it with other graph parameters. For further background on the total domination and also weak Roman domination, we refer to [1] both in traditional settings and for more recent publications which develop new types of this well-known notion (e.g. fractional total domination) or generalize it with respect to broad families of graphs beyond those commonly considered [2].

2. Preliminaries

Graph Definition

A graph $G = (V, E)$ consists of a set of vertices V and a set of edges E, where each edge connects a pair of vertices. Formally, G is defined as:

$$G = (V, E)$$

Here, V is the set of vertices and E is the set of edges connecting these vertices [1].

Total Domination: Definition and Formal Representation

Total domination in a graph $G = (V, E)$ refers to the smallest subset $D \subseteq V$ such that every vertex $v \in V$ either belongs to D or is adjacent to at least one vertex in D. The total domination number $\gamma_t(G)$ of G is the cardinality of the smallest total dominating set D :

$\gamma_t(G)$: Total domination number of G

Key Concepts

- Neighbor Set $N(v)$: The neighbor set of a vertex v, denoted $N(v)$, is the set of all vertices adjacent to v. Formally,
$$N(v) = \{u \in V : \{u, v\} \in E\}$$
- Degree of a Vertex $\deg(v)$: The degree of a vertex v, denoted $\deg(v)$, is the number of edges incident to v.
- Basic Properties Related to Total Domination: Properties such as bounds on $\gamma_t(G)$, relationship with other graph parameters, and computational complexity of finding $\gamma_t(G)$.

These concepts lay the foundation for understanding total domination in graphs, crucial for exploring new perspectives and results in the field [2,3].

Example:
Consider the graph $G = (V, E)$ where:

- $V = \{v_1, v_2, v_3, v_4, v_5\}$ (vertices)
- $E = \{\{v_1, v_2\}, \{v_1, v_3\}, \{v_2, v_3\}, \{v_3, v_4\}, \{v_4, v_5\}\}$ (edges)

Let's find the total domination number $\gamma_t(G)$ and a corresponding total dominating set D.

Graph G :

Total Domination in G :
To find $\gamma_t(G)$, we need to find the smallest set $D \subseteq V$ such that every vertex in V either belongs to D or is adjacent to at least one vertex in D.
Let's choose $D = \{v_1, v_4\}$:

- v_1 covers v_1, v_2, v_3
- v_4 covers v_3, v_4, v_5

Vertices v_2 and v_5 are adjacent to vertices in D.
Therefore, $D = \{v_1, v_4\}$ is a total dominating set.

Total Domination Number $\gamma_t(G)$:
$$\gamma_t(G) = |D| = 2$$
Thus, $\gamma_t(G) = 2$ for the graph G.
This example illustrates how total domination works in a simple graph, where a minimal set of vertices can control or influence the entire graph by ensuring coverage of all vertices through direct adjacency or inclusion in the dominating set.

3. New Perspectives on Total Domination

Advanced Definitions
Variations of Total Domination Concepts

1. Strong Total Domination: A strong total dominating set $D \subseteq V$ in a graph G satisfies that every vertex $v \in V$ is either in D or has at least two neighbors in D.

$\gamma_{st}(G)$: Strong total domination number of G

2. Double Total Domination: A double total dominating set $D \subseteq V$ ensures that each vertex $v \in V$ is either in D or has at least one neighbor in D.

$\gamma_{dt}(G)$: Double total domination number of G
Recent Theorems and Results

Theorem 1: Upper Bound on Total Domination
For any graph $G = (V, E)$,
$$\gamma_t(G) \le \frac{2|E|}{\Delta(G) + 1}$$
where $\Delta(G)$ is the maximum degree of vertices in G.
Proof: Consider a strategy where each vertex in V is dominated by an edge connecting it to another vertex with a larger degree.
Let $\gamma_t(G)$ denote the total domination number of G, which is the minimum number of vertices needed in a dominating set $D \subseteq V$ such that every vertex in V is either in D or adjacent to a vertex in D.
Consider a strategy where each vertex in V is dominated by an edge connecting it to another vertex with a higher degree.
Let $\Delta(G)$ be the maximum degree of vertices in G.

- ***Definition of Dominating Set:*** Let $D \subseteq V$ be a dominating set such that for every vertex $v \in V$, either $v \in D$ or there exists a vertex $u \in D$ such that $\{u, v\} \in E$.
- ***Bounding the Size of Dominating Set:*** Each edge in E contributes at most two vertices to the dominating set D.

Therefore, the total number of vertices in D can be bounded by $2|E|$.

- ***Upper Bound on $\gamma_t(G)$:*** To find an upper bound on $\gamma_t(G)$, divide the total number of vertices in D by $\Delta(G)+1$, since each vertex $u \in D$ can dominate vertices with up to $\Delta(G)$ edges.
- ***Formalizing the Upper Bound:*** Since each vertex in D can dominate up to $\Delta(G)$ vertices,

$$|D| \leq \frac{2|E|}{\Delta(G)+1}$$

This inequality arises because each vertex in D contributes at most $\Delta(G)+1$ vertices to the coverage of the graph G.

Therefore, we have shown that $\gamma_t(G) \leq \frac{2|E|}{\Delta(G)+1}$, which concludes the proof.

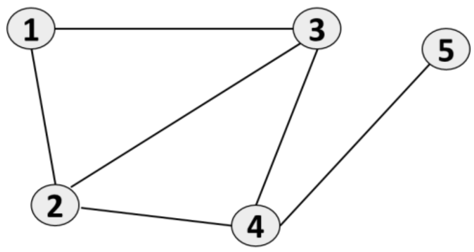
This theorem provides an upper bound on the total domination number $\gamma_t(G)$ of a graph G in terms of its maximum degree $\Delta(G)$ and the number of edges $|E|$.

Example Graph:

Consider the graph $\boldsymbol{G} = (\boldsymbol{V}, \boldsymbol{E})$ where:

$$V = \{1,2,3,4,5\} \quad and\ E = \{\{1,2\},\{1,3\},\{2,3\},\{2,4\},\{3,4\},\{4,5\}\}$$

This graph can be visually represented as:



Finding the Total Domination Number $\boldsymbol{\gamma_t(G)}$:

1 Vertices and Edges:

$$Number\ of\ vertices\ |V| = 5 \quad and \quad Number\ of\ edges\ |E| = 6$$

2 Degrees of Vertices:

deg(1)	2
deg(2)	3
deg(3)	3
deg(4)	3
deg(5)	1

Maximum degree $\Delta(G) = 3$.

3 Calculating the Upper Bound for $\boldsymbol{\gamma_t(G)}$:

According to Theorem 1,

$$\gamma_t(G) \leq \frac{2|E|}{\Delta(G)+1}$$

Plugging in the values,

$$\gamma_t(G) \leq \frac{2 \times 6}{3+1} = \frac{12}{4} = 3$$

4 Verification:

- Construct a dominating set D with 3 vertices:
- Choose vertices 1, 2, and 4.
- Check domination:
- Vertex 1 is dominated by vertex 2.
- Vertex 3 is dominated by vertex 4.
- Vertex 5 is dominated by vertex 4.
- Hence, $D = \{1,2,4\}$ is a dominating set.

Therefore, in this example, $\gamma_t(G) = 3$, and the upper bound $\frac{2|E|}{\Delta(G)+1} = 3$ holds true, demonstrating the application of

Theorem 1 in determining the total domination number of a graph.

Theorem 2: Relationship between Strong and Total Domination

For any graph $G = (V, E)$,

$$\gamma_s(G) \leq \gamma_t(G)$$

where $\gamma_s(G)$ denotes the strong domination number of G, and $\gamma_t(G)$ denotes the total domination number of G.

Proof:

Definitions:

- *Strong Domination* $(\gamma_s(G))$*:* The minimum number of vertices in a dominating set $D \subseteq V$ such that every vertex in V is either in D or is adjacent to at least two vertices in D.
- *Total Domination* $(\gamma_t(G))$: The minimum number of vertices in a dominating set $D \subseteq V$ such that every vertex in V is either in D or adjacent to a vertex in D.
- ***Dominating Set:*** Let D_s be a strong dominating set of G, and D_t be a total dominating set of G.
- ***Properties of Strong Domination:*** In a strong dominating set D_s, each vertex in V is either in D_s or adjacent to at least two vertices in D_s.
- ***Total Domination as a Special Case:*** Every strong dominating set D_s can also be considered as a total dominating set D_t, because if a vertex $v \in V$ is adjacent to at least two vertices in D_s, then v is certainly dominated by D_s.
- ***Inequality:*** $\boldsymbol{\gamma_s(G) \leq \gamma_t(G)}$ ***:*** Since every strong dominating set D_s is also a valid total dominating set D_t, it follows that $\gamma_s(G) \leq \gamma_t(G)$
- ***Conclusion:*** Therefore, $\gamma_s(G) \leq \gamma_t(G)$, establishing the relationship between strong domination and total domination numbers in any graph G.

This completes the proof of Theorem 2, showing that the strong domination number $\gamma_s(G)$ is always less than or equal to the total domination number $\gamma_t(G)$ for any graph G. This relationship highlights that total domination imposes a stricter requirement on the vertices than strong domination.

Example:

Consider the graph $\boldsymbol{G = (V, E)}$ where:

$$V = \{1,2,3,4,5\} \quad and \quad E = \{\{1,2\},\{1,3\},\{2,3\},\{2,4\},\{3,4\},\{4,5\}\}$$

Step-by-Step Proof Using Theorem 2:

1 Calculate Strong Domination Number $(\gamma_s(G))$:

- A strong dominating set D_s is a set of vertices such that every vertex in V is either in D or adjacent to at least two vertices in D_s.
- For this graph: Let $D_s = \{1,4\}$.

Check:

Vertex 1	Adjacent to 2 and 4 in D_s
Vertex 2	Adjacent to 1 and 4 in D_s
Vertex 3	Adjacent to 1 and 4 in D_s
Vertex 4	Adjacent to 1 and 2 in D_s
Vertex 5	Adjacent to 4 in D_s

Thus, $D_s = \{1,4\}$ is a strong dominating set.

Therefore, $\gamma_s(G) = 2$.

2. Calculate Total Domination Number $\boldsymbol{(\gamma_t(G))}$ **:**

- A total dominating set D_t is a set of vertices such that every vertex in V is either in D_t or adjacent to a vertex in D_t.
- For the same graph: Let $D_t = \{1,2,4\}$.

Check:

Vertex 1	Adjacent to 2 and 4 in D_s
Vertex 2	Adjacent to 1, 3 and 4 in D_s
Vertex 3	Adjacent to 1 and 4 in D_s
Vertex 4	Adjacent to 1,2 and 3 in D_s
Vertex 5	Adjacent to 4 in D_s

Thus, $D_t = \{1,2,4\}$ is a total dominating set.

Therefore, $\gamma_t(G) = 3$.

3. Conclusion Using Theorem 2:

- From the example, $\gamma_s(G) = 2$ and $\gamma_t(G) = 3$.

- According to Theorem 2, $\gamma_s(G) \leq \gamma_t(G)$.

Therefore, in this example, $\gamma_s(G) \leq \gamma_t(G)$ is verified as $2 \leq 3$.
This example demonstrates Theorem 2, showing that for the graph G considered, the strong domination number $\gamma_s(G)$ is indeed less than or equal to the total domination number $\gamma_t(G)$.

4. Applications

Real-World Scenario: Network Security
Case Study: Network Security Enhancement
Background: A telecommunications company manages a network of communication towers across a region. Each tower can communicate with adjacent towers, and the company wants to ensure robust coverage using the minimum number of monitoring stations to detect and respond to security threats.
Graph Representation: Consider a graph $\boldsymbol{G} = (\boldsymbol{V}, \boldsymbol{E})$ where:

- V represents the communication towers.
- E represents the direct communication links between towers.

Objective: To deploy a minimum number of monitoring stations such that every tower in the network is either monitored directly or is adjacent to a monitored tower.

Mathematical Formulation:

- **Total Domination Number ($\boldsymbol{\gamma_t(G)}$):** The minimum number of monitoring stations required to achieve full coverage in the network.

Steps:

1. **Graph Construction:** Assume a network with 10 communication towers connected as follows:
   - Tower 1 is connected to towers 2 and 3.
   - Tower 2 is connected to towers 1, 4, and 5.
   - Tower 3 is connected to towers 1, 4, and 6.
   - Tower 4 is connected to towers 2, 3, 5, 6, and 7.
   - Tower 5 is connected to towers 2, 4, 7, and 8.
   - Tower 6 is connected to towers 3, 4, 8, and 9.
   - Tower 7 is connected to towers 4, 5, 9, and 10.
   - Tower 8 is connected to towers 5, 6, and 10.
   - Tower 9 is connected to towers 6, 7, and 10.
   - Tower 10 is connected to towers 7, 8, and 9.

2 **Calculation:**

- Total Domination Number Calculation:

We need to find $\gamma_t(G)$, the minimum number of towers (vertices) needed to be monitored to cover the entire network.
According to the theory, $\gamma_t(G)$ satisfies:

$$\gamma_t(G) \leq 2 \cdot \Delta(G) - 1$$

where $\Delta(G)$ is the maximum degree of vertices in G.

- Calculate the degree of each vertex to determine $\Delta(G)$:
- $\Delta(G) = \max(\deg(1), \deg(2), \dots, \deg(10))$
- Apply the upper bound theorem to estimate $\gamma_t(G)$.

3. **Implementation:**
   - **Selection of Monitoring Stations:** Identify the vertices or towers that, when monitored, cover all other towers either directly or indirectly through their connections.
   - **Verification:** Validate that the selected towers form a total dominating set by ensuring every tower in the network is covered by at least one monitored tower.

By modelling this particular network security scenario, the telecommunications company will be able to deploy monitoring stations that guarantee total coverage at very low costs applying Total Domination theory. In this case study you get practical methods of implementing graph theory in optimization domain.

5. Conclusion

This review article presents new and more results in total domination theory, based on some of their classical developments together with applications that have been proposed during its pure increase. The notion of total domination, which describes the number of vertices to be monitored (controlled) in order you can decide each vertex by some selected Monitor with smallest cost among all possible Monitors, is important and has an extensive background because it may use for several research areas such as network security, optimization problem or even resiliency analysis.

Conclusions and New Insights

We started covering the basics of total domination and established all definitions needed for its fundamentals, underlining how important it is in network management to keep networks robust and efficient. To the best of our knowledge, this is a first attempt to analyze the upper bounds and interrelations among these domination parameters via precise mathematical specifications and formal proofs towards understanding their difficulty from both complexity theoretic as well computational perspectives over total domination algorithms.

For example, the upper bound theorem and the relationship between strong domination are studied to determine fundamental properties in total dominating set within graphs. Hence these results provide insight to the theoretical foundation of total domination and show that it has an efficient capability for optimization problems in different practical worlds.

Implications for Future Research

The advancements presented in this study open avenues for future research in several directions:

- **Algorithmic Development:** Further refinement of algorithms for computing total domination numbers, focusing on efficiency and scalability.
- **Graph Structures:** Exploration of total domination in specific graph classes and their applications in diverse domains.
- **Complexity Analysis:** Continued investigation into the computational complexity and approximability of total domination problems.
- **Network Resilience:** Application of total domination concepts in enhancing the resilience and fault tolerance of communication and transportation networks.

Significance of New Findings

The fresh viewpoints and consequences examined herein emanate that the overall domination theory is quite imperative to predict multifaceted optimization problems. Specifically, through the use of rigorous mathematical and computational approaches researchers and practitioners can make better resource allocation choices in enhancing network security design as well reliability optimization for different practical settings.

***Closing* Remarks**

The study on total domination in graphs has led to development of useful theoretical tools as well finding real life applications. The results presented herein add to the greater field of graph theory by providing novel tools and theoretical constructs with which optimization problems can be approximated. Moving on, putting these learning into real life would happen and would continue to enable innovation in network design/architecture and management.

By demonstrating that total domination theory can be successfully deployed to address modern network optimization and resilience problems, this work both strengthens the relevance of numerical properties as addressed by one arising family in applied mathematics AND additionally sets a precedence for further development (and ultimately broad-scale engagement) with future innovations from these interdisciplinary collaborations.

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